

On judicious bipartitions of directed graphs

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Taiwan, August 20, 2019

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Max-Cut problem

Let G be a graph. A **bipartition** of G , denoted by (V_1, V_2) , is a bipartition of $V(G)$ with $V(G) = V_1 \cup V_2$ and $V_1 \cap V_2 = \emptyset$.

If it satisfies $||V_1| - |V_2|| \leq 1$, then we call it a **bisection**.

The **size of (V_1, V_2)** , denoted by $e(V_1, V_2)$, is the number of edges with one end in V_1 and the other in V_2 .

The famous **Max-Cut problem** is to find a bipartition (V_1, V_2) of a given graph G that maximizes $e(V_1, V_2)$.

Brief analysis

Suppose that G has m edges and **chromatic number k** .

Let (V_1, \dots, V_k) be a k -coloring of G . Let G' be the **complete multigraph obtained from G by identifying each V_i into a single vertex**.

Then every bipartition of G' yields a bipartition of G with the same size.

Continued

By randomly partition $V(G)$ into two subsets of size $\lfloor k/2 \rfloor$ and $\lceil k/2 \rceil$, the expected number of edges between the two subsets is

$$\frac{\lfloor k/2 \rfloor \times \lceil k/2 \rceil}{\binom{k}{2}} m \geq \frac{(k^2 - 1)/4}{(k^2 - k)/2} m = \frac{m}{2} + \frac{m}{2k}.$$

Since $m \geq \binom{k}{2}$, we have

$$k \leq \sqrt{2m + \frac{1}{4}} + \frac{1}{2}.$$

Edwards's bound

This implies that G has a bipartition (V_1, V_2) such that

$$e(V_1, V_2) \geq \frac{m}{2} + \frac{1}{4} \left(\sqrt{2m + \frac{1}{4}} - \frac{1}{2} \right),$$

and this bound is best possible for K_{2n+1} .



C.S. Edwards, Some extremal properties of bipartite graphs, *Canad. J. Math.* **3** (1973) 475–485.



C.S. Edwards, An improved lower bound for the number of edges in a largest bipartite subgraph, *Proc. 2nd Czechoslovak Symposium on Graph Theory* (1975) 167–181.

Triangle-free graphs

An natural direction of the Max-Cut problem is to bound the Max-Cut of graphs without a specific subgraph H . The Max-Cut problem of **triangle free graphs** was posed by Erdős.

Theorem

Every triangle free graph G with m edges admits a bipartition (V_1, V_2) such that

$$e(V_1, V_2) \geq m/2 + cm^{2/3},$$

for some $c > 0$.



P. Erdős, Problems and results in graph theory and combinatorial analysis, In Graph Theory and Related Topics (Proc. Conf. Waterloo, 1977), Academic Press, New York (1979) 153–163.

Shearer's bound

Let d_1, \dots, d_n be the degree sequence of a triangle free graph G with n vertices and m edges. Shearer showed that G has a bipartition of size at least

$$m/2 + \frac{1}{8\sqrt{2}} \sum_{i=1}^n \sqrt{d_i}.$$

It follows as a corollary that G has a bipartition of size at least

$$m/2 + cm^{3/4},$$

for some $c > 0$.



J. Shearer, A note on bipartite subgraphs of triangle-free graphs, Random Struct. Alg. 3 (1992) 223–226.

Alon's bound

In 1996, Alon proved

Theorem

Every triangle free graph G with m edges admits a bipartition (V_1, V_2) such that

$$e(V_1, V_2) = m/2 + \Theta(m^{4/5}).$$



N. Alon, Bipartite subgraphs, *Combinatorica* 16 (1996) 301–311.

Large girth

For $r \geq 3$, let C_r denote a cycle of length r . Alon, Bollobás, Krivelevich and Sudakov considered the maximum cut of graphs G with large **girth**, and proved that **every graph G with m edges and girth at least r admits a bipartition (V_1, V_2) such that**

$$e(V_1, V_2) \geq m/2 + c(r)m^{\frac{r}{r+1}}$$

for some $c(r) > 0$.



N. Alon, B. Bollobás, M. Krivelevich, B. Sudakov, Maximum cuts and judicious partitions in graphs without short cycles, J. Combin. Theory Ser. B 88 (2003) 329–346.

Even cycle

Alon, Krivelevich and Sudakov considered graphs without even cycles and proved that

Theorem

For every *odd* $r \geq 5$, and a C_{r-1} -free graph G with m edges, there is a positive constant $c(r)$ such that G has a bipartition of size at least

$$m/2 + c(r)m^{\frac{r}{r+1}}.$$



N. Alon, M. Krivelevich and B. Sudakov, Maxcut in H -free graphs, *Combin. Probab. Comput.* 14 (2005), 629–647.

Conjecture

In the same paper, they conjectured that the lower bound holds for graphs without odd cycles.

Conjecture

For every *even* integer $r \geq 4$, and a C_{r-1} -free graph G with m edges, there is a positive constant $c(r)$ such that G has a bipartition of size at least

$$m/2 + c(r)m^{\frac{r}{r+1}}.$$

So far, this conjecture was confirmed for $r = 4$.

Our result

We considered the conjecture and established the following theorem.

Theorem (H., Zeng, 2018, ARS Math. Contemp.)

For every even integer $r > 4$, and a C_{r-1} -free graph G with m edges, there is a positive constant $c(r)$ such that G has a bipartition of size at least

$$m/2 + c(r)(m^r \log^4 m)^{\frac{1}{r+2}}.$$

5-Cycles free

For C_5 -free graphs, we have

Theorem (H., Zeng, 2018, ARS Math. Contemp.)

For each $s \geq 2$, let G be a $\{K_{2,s}, C_5\}$ -free graph with m edges. Then G has a bipartition of size at least

$$m/2 + c(s)m^{6/7}$$

for a positive constant $c(s)$.

General graph

It is noted that for every fixed graph H there exist positive constants $\epsilon = \epsilon(H)$ and $c = c(H)$ satisfying that:

if a H -free graph G has m edges, then G has a bipartition of size at least

$$m/2 + cm^{1/2+\epsilon}.$$

A conjecture

Alon, Bollobás, Krivelevich and Sudakov posed the following conjecture.

Conjecture

Let H be a fixed graph, and let G be a graph with m edges. If G is H -free, then G has a bipartition of size at least

$$\frac{m}{2} + \Omega(m^{3/4+\epsilon}).$$

for some $\epsilon > 0$.



N. Alon, B. Bollobás, M. Krivelevich, B. Sudakov, Maximum cuts and judicious partitions in graphs without short cycles, J. Combin. Theory Ser. B 88 (2003) 329–346.

Clearly, it suffices to prove this conjecture for complete graphs H . We considered K_{k+1} -free graph and proved that

Theorem (H., Zeng, 2017, Bull. Aust. Math. Soc.)

For any fixed integer $k \geq 2$, every K_{k+1} -free graph with m edges admits a bipartition of size at least

$$\frac{m}{2} + c(k)m^{\frac{k}{2k-1}} \left(\frac{\log^2 m}{\log \log m} \right)^{\frac{k-1}{2k-1}}.$$

for some $c(k) > 0$.

Max-Bisection

The **Max-Bisection problem** is that given a graph G , find a **bisection** (V_1, V_2) of G that maximizes $e(V_1, V_2)$.

Compared to bipartitions, bisections are much more complicated to analyze.

For example, Edwards' bound implicitly implies that a **connected graph G with n vertices and m edges admits a bipartition (V_1, V_2) with**

$$e(V_1, V_2) \geq m/2 + (n - 1)/4.$$

Unfortunately, this result cannot transfer directly to bisections.

Example

Let $K_{k,n-k}$ be the complete bipartite graph with $k \leq n/2$.

It is easy to see that, the size of the maximum bisection of $K_{k,n-k}$ is $m/2 + k^2/2$ if n is even, and it is $m/2 + (k^2 + k)/2$ if n is odd, where $m = k(n - k)$.

So, for bisections, we cannot get a bound greater than $(m + 1)/2$ for every graph G with m edges.

Maximum matching

Let G be a graph with m edges, and let M be a maximum matching of G .

Then an easy analysis, given by Xu, Yan and Yu, shows that G has a bisection of size at least

$$\frac{m + |M|}{2}.$$



B. Xu, J. Yan, X. Yu, A note on balanced bipartitions, Discrete Math. 310 (2010) 2613–2617.

Tight graph

A connected graph T is **tight** if

- for every vertex $v \in V(T)$, $T - v$ contains a perfect matching, and
- for every vertex $v \in V(T)$ and every perfect matching M of $T - v$, no edge in M has exactly one end adjacent to v .

Lu, Wang and Yu showed that **a connected graph G is tight iff every block of G is an odd clique.**



C. Lu, K. Wang, X. Yu, On tight components and anti-tight components, *Graphs and Combinatorics* 31 (2015) 2293–2297.

Tight bound

Lee, Loh and Sudakov gave a lower bound of Max-Bisection of graphs with respect to the tight components and maximum degree by showing

Theorem

Let G be a graph with m edges and maximum degree Δ . If G has τ *tight components*, then there is a bisection of size at least

$$\frac{m}{2} + \frac{n - \max\{\tau, \Delta - 1\}}{4}.$$

The vertex-disjoint copies of a triangle and the star $K_{1,n-1}$ show that the theorem is tight in both parameters τ and Δ .



C. Lee, P. Loh, B. Sudakov, Bisections of graphs, J. Combin. Theory Ser. B 103 (2013) 590–629.

Our result

The following result showed a similar bound on bisections of graphs without short cycles.

Theorem (Fan, H., Yu, 2018, CPC)

Let G be a graph with n vertices, m edges and minimum degree $\delta \geq 2$, and without 4-cycles. If $\delta(G)$ is even, then G has a bisection of size at least

$$\frac{m}{2} + \frac{n-1}{4} - \frac{|M|}{2\delta},$$

where M is a maximum matching in G . Moreover, if the girth $g(G) \geq 5$, then G has a bisection of size at least

$$\frac{m}{2} + \frac{n-1}{4}.$$

Generalized results (1)

Note that the **distance** between two triangles is the length of a shortest path between their vertices.

Jin and Xu generalized the above result by giving the following theorem.

Theorem

*Let $l \geq 2$ be an integer, and let G be a connected graph with minimum degree at least 2 and **without $K_{2,l}$** . Then, G admits a bisection (V_1, V_2) with $e(V_1, V_2) \geq m/2 + (n - l + 1)/4$ if either G is **Eulerian**, or G has **no triangles at distance 2**.*



J. Jin and B. Xu, Bisection of graphs without $K_{2,l}$, Discrete Appl. Math. 259 (2019) 112–118.

Generalized results (2)

A (l, r) -fan is a graph on $(r - 1)l + 1$ vertices consisting of l cliques of order r that intersect in exactly one common vertex.

Theorem (H., Yan, 2019, DM)

Let $l \geq 2$ be an integer, and let G be a connected graph with m edges and minimum degree at least 2. If G contains neither $K_{2,l}$ nor $(l, 4)$ -fan, then G admits a bisection (V_1, V_2) with

$$e(V_1, V_2) \geq \frac{m}{2} + \frac{n - l + 1}{4}.$$

Corollary

When $l = 2$, we have the following corollary.

Corollary

*Every connected graph G with m edges and minimum degree at least 2, and **without 4-cycles** admits a bisection of size at least*

$$\frac{m}{2} + \frac{n-1}{4}.$$

Bondy-Simonovits Theorem

The following result of Bondy and Simonovits gives the maximum number of edges in graphs without cycles of a given even length.

Theorem

Let $l \geq 2$ be an integer and let G be a graph with n vertices. If G contains no cycle of length $2l$, then the number of edges in G is at most $100ln^{1+1/l}$.



A. Bondy, M. Simonovits, Cycles of even length in graphs, J. Combin. Theory Ser. B 16 (1974) 97–105.

New bound

Combining Corollary 8 and Theorem 9, we have

Theorem

Let G be a graph with m edges and minimum degree $\delta \geq 2$. If G contains no 4-cycle, then G has a bisection of size at least

$$m/2 + cm^{2/3},$$

for some constant $c > 0$.

Definition

Given a digraph D , a digraph version of the Max-Cut problem is to find a partition (V_1, V_2) of D that **maximizes**

$$\min\{e(V_1, V_2), e(V_2, V_1)\}.$$

Here, $e(V_1, V_2)$ is the number of arcs xy with $x \in V_1$ and $y \in V_2$.

Such a problem are called **Judicious Partition Problem** by Bollobás and Scott.



B. Bollobás, A.D. Scott, Problems and results on judicious partitions, *Random Struct. Alg.* 21 (2002) 414–430.

Scott's problem

Scott posed the following natural problem:

What is the maximum constant c_d such that every digraph D with m arcs and minimum outdegree d admits a bipartition (V_1, V_2) satisfying that

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq c_d m?$$



A.D. Scott, Judicious partitions and related problems, *Surveys in Combinatorics* 327 (2005) 95–117.

$d=1$

For $d = 1$, consider the graph $K_{1,n-1}$ and add a single edge inside the part of size $n - 1$. This graph can be oriented easily so that the minimum outdegree is 1 and

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \leq 1$$

for every partition (V_1, V_2) of such a graph. Hence,

$$c_1 = 0.$$

Lee-Loh-Sudakov Conjecture

Lee, Loh and Sudakov initiated the study of this problem and conjectured that

Conjecture

Let d be an integer satisfying $d \geq 2$. Every digraph D with m arcs and minimum outdegree at least d admits a bipartition (V_1, V_2) for which

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq \left(\frac{d-1}{2(2d-1)} + o(1)\right)m.$$



C. Lee, P. Loh, B. Sudakov, Judicious partitions of directed graphs, *Random Struct. Alg.* 48 (2016) 147–170.

Extremal graph

The bound in Conjecture 3 is **asymptotically best possible** by considering the following extremal graph:

First orient the arcs of the complete graph K_{2d-1} along an Eulerian circuit.

Then consider the directed graph where we take **k vertex disjoint copies of K_{2d-1}** oriented as above, and a single vertex disjoint copy of K_{2d+1} oriented in a similar manner.

Fix a vertex v_0 of K_{2d+1} , and **add arcs** so that all the vertices belonging to the copies of K_{2d-1} are in-neighbors of v_0 .

Continued

Clearly, the resulting graph has **minimum outdegree** d ,
minimum degree $2d - 1$, **minimum semidegree** $d - 1$.

Progress

In the same paper, they confirmed the conjecture when $d = 2, 3$.

On the other hand, they noted that the methods they used turn out to be too limited in strength to cover the cases $d \geq 4$.



C. Lee, P. Loh, B. Sudakov, Judicious partitions of directed graphs, *Random Struct. Alg.* 48 (2016) 147–170.

Dense case

For digraphs with bounded maximum degree or large number of arcs, we can get a better bound.

Theorem (H., Wu, Yan, EJC, 2017)

Let D be a digraph with n vertices and m arcs. For every $\epsilon > 0$, if $m \geq 16n/\epsilon^2$ or the maximum degree Δ of D is at most $\epsilon^2 m/8$, then D admits a bisection (V_1, V_2) such that

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq \frac{1}{4}m - \epsilon m.$$

Standard approach

A standard approach to find a “good” bipartition of a digraph D is to first **partition $V(D)$ into X, Y** , where X consists of certain **high degree** vertices.

Then, partition X into **X_1 and X_2** with certain property;

Finally, apply a **randomized algorithm** to distribute the vertices in Y .



B. Bollobás, A.D. Scott, Judicious partitions of hypergraphs, J. Combin. Theory Ser. A 78 (1997) 15–31.

Difficulty

The main step in this approach is to deal with **the arcs between X and Y** .

The condition that **“the minimum outdegree of D is at least d ”** yields that the number of arcs from Y to X can be bounded.

However, if we bound the number of arcs from X to Y , we need to know the **indegree of vertices in Y** .

Semidegree

Our first result shows that Conjecture 3 holds under the natural (additional) assumption that **the minimum indegree of D is at least d** .

Theorem (H., Ma, Yu, Zhang, 2019, SCM)

Let $d \geq 2$ be an integer. Every digraph D with m arcs and **minimum semidegree** at least d admits a bipartition (V_1, V_2) for which

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq \left(\frac{d-1}{2(2d-1)} + o(1)\right)m.$$

Oriented graph

Note that an **oriented graph** is an orientation of a simple graph. If we focus on oriented graphs with minimum semidegree d , then we can give a better bound.

Theorem (H., Wu, JCTB, 2018)

Let d be an integer with $d \geq 21$. Every **oriented** graph D with m arcs and **minimum semidegree d** admits a bipartition (V_1, V_2) such that

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq \left(\frac{d}{2(2d+1)} + o(1)\right)m.$$

Special digraphs

Let $\overrightarrow{K_{d,2}}$ be the digraph obtained by orient each edge of a bipartite graph $K_{2,d}$ from the part of size d to the other part.

Theorem (H., L., Wu, 2019+)

Let $d \geq 4$ be an integer and D is a digraph with m arcs and minimum outdegree at least d . If D does not contain $\overrightarrow{K_{d,2}}$, then there is a bipartition (V_1, V_2) such that

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq \left(\frac{d-1}{2(2d-1)} + o(1)\right)m.$$

Underlying graph

The **underlying graph** of a digraph is obtained by ignoring arc orientations and removing redundant parallel arcs when arcs in both directions appear between pairs of vertices.

Without 4-cycles

For digraphs whose **underlying graph does not have 4-cycles**, we have

Theorem (H., Wu, Yan, EJC, 2017)

*Let D be a digraph with m arcs and **minimum outdegree at least 2**, and let G be its underlying graph. If G **does not contain 4-cycles**, then D admits a bisection (V_1, V_2) such that*

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq \left(\frac{1}{4} + o(1)\right)m.$$

Conclusion (1)

The Max-Cut problem is a fundamental discrete optimization problem and have been studied widely. We give the lower bound on Max-Cut of graphs without specific structure.

There are litter results on Max-Bisections of graphs. Even for C_4 -free graphs G , we do not know the maximum constant c such that the following holds:

If G has m edges and minimum degree $\delta \geq 2$, then G has a bisection of size at least

$$m/2 + \Omega(m^c).$$

Conclusion (2)

We conclude our discussion with the following question:

Question (H., Wu, JCTB, 2018)

Is it true that for every integer $d \geq 1$, every digraph D with m arcs and **minimum semidegree d** admits a bipartition $V(D) = V_1 \cup V_2$ such that

$$\min\{e(V_1, V_2), e(V_2, V_1)\} \geq \left(\frac{d}{2(2d+1)} + o(1)\right)m?$$

As evidence, we have showed that it is true for digraphs with minimum semidegree at most 3.

Thank you for your attention!