Disjoint cycles in digraphs

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- 1. Notation
- 2. Introduction and Main Theorem
- 3. Sketch of the Main Theorem
- 4. Related Results and Open Problems

1. Notation

• A graph G = (V, E): V vertex set, E edge set.

Let n be the order of G for simplicity, i.e. n=|V|. $\delta(G):$ the minimum degree of G.

A digraph D = (V, A): V vertex set, A arc set.
 δ⁺(D): the minimum out-degree,
 δ⁻(D): the minimum in-degree of D.

• The semi-degree of D is $\delta^0(D) = \min\{\delta^+(D), \delta^-(D)\}.$

In a digraph: a cycle is always **directed**.

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- In a digraph: a cycle is always **directed**.

- A **tournament** is a digraph *T* such that for any two distinct vertices *x* and *y*, exactly one of the ordered pairs (*x*, *y*) and (*y*, *x*) is an arc of *T*.
- A set of subgraphs of G or D is said to be vertex-disjoint(briefly disjoint) if no two of them have any common vertex in G or D.
- A 2-factor of a graph G is a spanning subgraph of G such that each component is a cycle.
 A hamiltonian cycle is a 2-factor with exactly one

component.

■ A cycle factor of a digraph *D* is a spanning subgraph of *D* such that each component is a cycle in *D*.

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2. Introduction and Main Theorem

For a graph G, if $e(G) \ge n$, or $\delta(G) \ge 2$, then G contains a cycle.

How about two disjoint cycles?



P. Erdős, L. Pósa, 1962

For every graph G, if $n \ge 6$ and $e(G) \ge 3n - 6$, then G has 2 disjoint cycles or isomorphic to $K_3 + (n - 3)K_1$.

2. Introduction and Main Theorem

Given an integer $k \ge 1$, how about k disjoint cycles?

K. Corrádi and A. Hajnal, 1963

For any positive integer k and any graph G, if $n \ge 3k$ and $\delta(G) \ge 2k$, then G has k disjoint cycles.

Sharpness of degree condition is given by $G = K_{2k-1} + mK_1$.



m independent vertices

$$G = K_{2k-1} + mK_1$$

└─2. Introduction and Main Theorem





G. A. Dirac, 1963

For any positive integer k and any graph G, if $n \ge 3k$ and $\delta(G) \ge (n+k)/2$, then G contains k disjoint triangles.

 The papers in this topic can be found in [Degree Conditions for the Existence of Vertex-Disjoint Cycles and Paths: A Survey, Graphs and Combinatorics (2018) 34:1 - 83.] For a digraph D, if $\delta^+(D) \ge 1$, then D contains a cycle.

Given a positive integer k and a digraph D, what is the degree condition for k disjoint cycles in a digraph?

Conjecture 2.1 [Bermond, Thomassen, J. Graph Theory 5 (1) (1981) 1-43.]

For $k \ge 1$ and any digraph D, if $\delta^+(D) \ge 2k - 1$, then D contains k disjoint cycles.

- Thomassen proved the case k = 2 in 1983 (Combinatorica);
- Lichiardopol et al. proved the case k = 3 in 2009 (SIAM Discrete Math.);
- Conjecture 2.1 remains open for $k \ge 4$.

Thomassen further proposed the conjecture on disjoint cycles with the same length.

Conjecture 2.2 [Thomassen, Combinatorica 3 (3-4) (1983) 393-396.]

For each natural k and any digraph D, there exists g(k), if $\delta^+(D) \geq g(k)$, then D contains k pairwise disjoint cycles of the same length.

In 1996, Alon disproved the conjecture.

Theorem 2.3 [Alon, J. Combin. Theory Ser. B 68 (2) (1996) 167-178.]

For every integer r, there exists a digraph D such that $\delta^+(D) = r$, but D contains no two edge-disjoint cycles of the same length (and hence, of course, no two vertex-disjoint cycles of the same length).

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Although Conjecture 2.2 is not true for general digraphs, Lichiardopol believed that it is correct for tournaments. He raised the following conjecture.

Conjecture 2.4 [Lichiardopol, Discrete Math. 310 (19) (2010) 2567-2570.]

For every tournament T, if $\delta^+(T) \ge (q-1)k - 1$, then T contains k disjoint cycles of length q, where $q \ge 3$ and $k \ge 1$.

In the same paper, Lichiardopol himself proved that if both the minimum out-degree and in-degree are at least (q-1)k-1, i.e. $\delta^0(T) \ge (q-1)k-1$, then T contains k disjoint cycles of length q.

- In 2013, Jensen, Bessy and Thomassé proved Conjecture 2.4 for the case q = 3 (A special case of Bermond-Thomassen conjecture).
- We confirm Lichiardopol's Conjecture for $q \ge 4$.

Main Theorem 1 [F. Ma, Y. Wang and Y, 2019+]

Conjecture 2.4 is correct.

That is, for any integer $q \ge 4$, $k \ge 1$ and every tournament T, if $\delta^+(T) \ge (q-1)k-1$, then it contains k disjoint cycles of length q.

We also improved Lichiardopol's theorem on semi-degree condition significantly by proving the following theorem.

Main Theorem 2 [F. Ma and Y, Applied Mathematics and Computation 347 (2019) 162-168.]

For integers $k \ge 1$ and $q \ge 4$, every tournament T with $\delta^0(T) \ge (q-1)k - 1$ contains f(q)k - 2q disjoint cycles of length q, where $f(q) = \frac{6q^2 - 16q + 10}{3q^2 - 3q - 4}$.

Note

- f(q) > 1, when $q \ge 4$
- f(q) tends to 2, when q tends to infinity.

- └─3. Sketch of Main Theorem 1
 - 3. Sketch of Main Theorem 1

Property 1. Every tournament has a hamiltonian path.

Property 2. Every strong tournament is vertex pancyclic.

- We prove Main Theorem by induction on k.
- When k = 1, then $\delta^+(T) \ge q 2$. Let $P = (v_n \cdots v_1)$ be a hamiltonian path of T.

Consider v_1 , since $d^+(v_1) \ge q - 2$ and all its out-neighbours are on P, there exists v_i with $i \ge q$ satisfying $v_1v_i \in A$.

■ So there is a cycle $v_i v_{i-1} \cdots v_1 v_i$ of length at least q. By property of pancyclicity, T contains a q-cycle (a cycle of length q).

- Suppose that Main Theorem is correct for k-1, i.e. if $\delta^+(T) \ge (q-1)(k-1)-1$, then T contains k-1 disjoint q-cycles.
- Now we consider the case k. Since in this case $\delta^+(T) \ge (q-1)k - 1 > (q-1)(k-1) - 1$, T contains k - 1 disjoint q-cycles.
- Using the following three auxiliary theorems, we may show that T contains k disjoint q-cycles,

Theorem 3.1

Let $q \ge 9$ and $k \le q+1$. For every collection \mathcal{F} of k-1 disjoint q-cycles of T, there exists a collection of k disjoint q-cycles in T which intersects $T \setminus \mathcal{F}$ on at most 3q vertices.

Theorem 3.2

Let $q \geq 10$ and $k \geq 3.3098\sqrt{q}$. For every collection \mathcal{F} of k-1 disjoint q-cycles of T, there exists a collection of k disjoint q-cycles in T which intersects $T \setminus \mathcal{F}$ on at most 3q vertices.

• Since $3.3098\sqrt{q} \le q+2$ when $q \ge 8$, Theorems 3.1 and 3.2 imply Main Theorem for $q \ge 10$.

Theorem 3.3

When $4 \le q \le 9$, we can refine Theorem 3.2 with " $|\mathcal{F} \cap (T \setminus \mathcal{F})| \le 3q - 6$ ". That is, For every collection \mathcal{F} of k - 1 disjoint q-cycles of T, there exists a collection of k disjoint q-cycles in T which intersects $T \setminus \mathcal{F}$

on at most 3q - 6 vertices.

Let $\mathcal{F} = \{C_1, \dots, C_{k-1}\}$ be a collection of k-1 disjoint q-cycles, $P = (u_r \cdots u_2 u_1)$ be a hamiltonian path of $T \setminus \mathcal{F}$.

We design the following pseudo algorithm.

step 0. Set $\mathcal{F} := \{C_1, \ldots, C_{k-1}\}$, $P := (u_r \cdots u_2 u_1)$ and h := 0. Let $C = (u_{j-1}u_{j-2} \cdots u_1 u_{j-1})$ be the longest cycle of T[V(P)] through vertex u_1 , if it exists.

step 1. If C does not exist, then we construct a collection \mathcal{F}' of k-1 disjoint q-cycles and a new hamiltonian path $P' = (v_r \cdots v_2 v_1)$ in $T \setminus \mathcal{F}'$ such that $|V(\mathcal{F}') \cap V(P)| \leq 2$ and $T \setminus \mathcal{F}'$ contains a cycle C' through vertex v_1 . Set $\mathcal{F} := \mathcal{F}', P := P', C := C'$ and h := h + 1. Otherwise, go to step 2.

└─3. Sketch of Main Theorem 1

3. Sketch of Theorem 3.1

step 2. Construct a collection \mathcal{F}' of k-1 disjoint q-cycles and a new hamiltonian path $P' = (v_r \cdots v_2 v_1)$ in $T \setminus \mathcal{F}'$ such that $|V(\mathcal{F}') \cap V((P))| \leq 2$ and $T \setminus \mathcal{F}'$ contains a cycle C' with $|C'| \geq |C|+1$ through vertex v_1 . Set $\mathcal{F} := \mathcal{F}'$, P := P', C := C' and h := h + 1.



step 3. If $|C| \le q - 1$, then do step 2 recursively until $|C| \ge q$. Otherwise, output $\mathcal{F} \cup \{C\}$ and h.

- After step 3, |C| ≥ q, by pancyclicity, there is a q-cycle C'. This cycle and *F* form a collection of k disjoint q-cycles.
- Now we prove that the vertex number of $|(V(\mathcal{F} \cup \{C'\})) \cap V(P)|$ is at most 3q by using the algorithm.
- At each iteration of step 1 and step 2, we add at most two vertices outside $\{C_1, \ldots, C_{k-1}\}$ into \mathcal{F} . It follows by $h \leq q-2$ that

$$|V(\mathcal{F}) \cap V((T \setminus \{C_1, \dots, C_{k-1}\}))| \le 2h \le 2q - 4.$$

Therefore,

$$|(\mathcal{F} \cup \{C'\}) \cap (T \setminus \{C_1, \dots, C_{k-1}\})| \le 2q - 4 + q = 3q - 4.$$

Thus Theorem 3.1 holds.

Let $\mathcal{F} = \{C_1, \ldots, C_{k-1}\}$ be a collection of k-1 disjoint q-cycles and let $P = (u_r \cdots u_2 u_1)$ be a hamiltonian path of $T \setminus \mathcal{F}$. Partition P by letting $U_1 = \{u_1, \ldots, u_{q+1}\}$, $S = \{u_{q+2}, \ldots, u_{4q-5}\}$ and $U_2 = V(P) \setminus (U_1 \cup S)$. That is, U_1 is the set of the last q + 1 vertices on P and S is the last 3q - 6 vertices on $P \setminus U_1$.



Fig.2 Partition of P

└─3. Sketch of Main Theorem 1

3. Sketch of Theorem 3.2

Denote by \mathcal{I} the set of *q*-cycles that receive at least q^2 arcs each from U_1 , by \mathcal{O} the set of *q*-cycles that send at least 6q-1 arcs each to U_2 and $\mathcal{R} = \mathcal{F} \setminus (\mathcal{I} \cup \mathcal{O})$. Furthermore, *i*, *o* and *r*, respectively, denote the size of \mathcal{I}, \mathcal{O} and \mathcal{R} .



Fig.3 Partition of F

└─3. Sketch of Main Theorem 1

3. Sketch of Theorem 3.2

Now we estimate the lower and upper bound of the number of arcs leaving from $\mathcal{F}\setminus\mathcal{O}=\mathcal{I}\cup\mathcal{R}.$



Fig.3 Partition of F

First, since $\delta^+(T) \ge (q-1)k - 1$,

$$d^{+}(\mathcal{I} \cup \mathcal{R}) \ge q(i+r)((q-1)k-1) - \frac{1}{2}q(i+r)(q(i+r)-1).$$

On the other hand, we bound the number of arcs from \mathcal{I} to \mathcal{O} and \mathcal{R} to \mathcal{O} , from \mathcal{I} to U_2 and \mathcal{R} to U_2 , from $\mathcal{I} \cup \mathcal{R}$ to S and U_1 .

 $d^{+}(\mathcal{I} \cup \mathcal{R}) \le (q^{2} - q + 2)io + q^{2}ro + (3q - 1)i + (6q - 2)r + (3q - 6)q(i + r)$

$$+q(i+r)(q+1)-\alpha+(q^2-q-2)o,$$
 where $\alpha=(q+1)((q-1)k-1)-\frac{1}{2}(q+1)q-\frac{1}{2}(q-2)(q-3).$

So we get

$$ao^2 + bo + c < 0, \tag{1}$$

where

$$a = \frac{1}{2}q^2 + q - 2, \ b = ((q - q^2)k) + (3q^2 + \frac{13}{2}q - 9)$$

and

$$c = (\frac{1}{2}q^2 - q)k^2 + (1 - \frac{1}{2}q - 3q^2)k + (\frac{5}{2}q^2 + \frac{11}{2}q - 25).$$

Obviously, $a = \frac{1}{2}q^2 + q - 2 > 0$. Inequality (1) admits solution for o only if

$$\Delta = (-2q^3 + 9q^2 - 8q)k^2 + (6q^3 + 7q^2 - 26q + 8)k + 4q^4 + 18q^3 + \frac{145}{4}q^2 + 27q - 119 > 0.$$
(2)

Note that (2) is a quadratic inequality for k. Since $-2q^3 + 9q^2 - 8q < 0$, the inequality (2) has a solution only if

$$k < g(q),$$

where

$$g(q) = \frac{6q^3 + 7q^2 - 26q + 8 + f(q)}{4q^3 - 18q^2 + 16q},$$

and

$$f(q) = \sqrt{32q^7 + 36q^6 - 146q^5 - 776q^4 - 1032q^3 + 5936q^2 - 4224q + 64}.$$

It follows by $q \ge 10$ that $k < 3.3098\sqrt{q}$, i.e. the inequality (2) has a solution only if $k < 3.3098\sqrt{q}$. This contradicts $k \ge 3.3098\sqrt{q}$. So Theorem 3.2 is proved. \Box

4. Refine Theorem 3.2 for $4 \le q \le 9$ Let $U_1 = \{u_1, \ldots, u_q\}$, $S = \{u_{q+1}, \ldots, u_{4q-10}\}$ and $U_2 = V(P) \setminus (U_1 \cup S)$ (If |P| < 4q - 10, then let $U_1 = \{u_1, \ldots, u_q\}$, $S = V(P) \setminus U_1$). Define

$$\mathcal{I} = \{ C \in \mathcal{F} \, | \, d^+(U_1, C) \ge q(q-1) + 1 \},\$$

$$\mathcal{O} = \{ C \in \mathcal{F} \, | \, d^+(C, U_2) \ge 6q - 13 \} \text{ and } \mathcal{R} = \mathcal{F} \setminus (\mathcal{I} \cup \mathcal{O}).$$

• Let C be a q-cycle. If there is a q-matching M from U_1 to C, then $d^+(C,S)\leq \frac{9}{4}q^2-\frac{29}{4}q.$

Similarly, estimate the lower and upper bound of $d^+(\mathcal{F} \setminus \mathcal{O})$. We get all the possible cases: if q = 4, then $1 \le k \le 4$; if $9 \ge q \ge 5$, then $1 \le k \le 5$.

From the following statement, we finish the proof of the case $4 \le q \le 9$.

• Let k be an integer with $k \leq 5$. If there exist two cycles $C_1, C_2 \in \mathcal{F}$ such that $d^+(\{u_1, u_2\}, C_i) \geq 2q - 1$ for i = 1, 2, then we can extend \mathcal{F} .

- 4. Related Results and Open Problems -Cycle Factor in Digraphs
- For a subset $W \subseteq V(D)$, define

$$\delta^+(W) = \min \{d^+_D(v) : v \in W\},\$$

$$\delta^-(W)= \min \ \{d^-_D(v): v\in W\}.$$

• The minimum semi-degree of W in D: $\delta^0(W) = \min \{\delta^+(W), \delta^-(W)\}.$

Theorem 4.1 [Y. Wang and Y, 2019+]

Suppose that D is a digraph with order n and $W \subseteq V(D)$. If $\delta^0(W) \ge (3n-3)/4$, then for any k positive integers n_1, \ldots, n_k with $n_i \ge 2$ for all i and $\sum_{i=1}^k n_i \le |W|$, D contains k disjoint cycles C_1, \ldots, C_k such that $|V(C_i) \cap W| = n_i$ for each i.

4. Related Results and Open Problems

- A directed version of the Aigner-Brandt Theorem, when W = V(D) [J. Lond. Math. Soc. 1 (1993) 39-51]: If $\delta(G) \ge (2n-1)/3$, then G contains k disjoint cycles of length n_1, \ldots, n_k , respectively. $(n \ge \sum_{i=1}^k n_i \text{ and } n_i \ge 3 \text{ for all } i)$
- Sharp (in some sense)



 $D_1: U = X = K_{4k-1}^*, Y = Z = K_{4k}^*$

 $D_2: X = K^*_{2k-1}$, Y is an independent vertex set of order k+1.

4. Related Results and Open Problems

Conjecture 4.2 [Y. Wang and Y, 2019+]

The minimum semi-degree in Theorem 4.1 can be improved to 2n/3 when $n_i \ge 3$ for all i.

The degree condition is best possible by D_2 .

It is supported by the following conjecture.

Conjecture 4.3 [Czygrinow, Kierstead and Molla, Eur. J. Combin. 42 (2013) 1-14] If n = 3k and $\delta^0(D) \ge 2k$, then D contains k disjoint \triangle s.

4. Related Results and Open Problems

Remark for Theorem 4.1.

 \blacksquare When $k=1,\,\delta^0(W)\geq (3n-3)/4 \Longrightarrow \delta^0(W)\geq \frac{n}{2}$

• Let
$$\lambda = \sum_{i=1}^{k} n_i$$
.
If $n \ge 2\lambda$, then $\delta^0(W) \ge (3n-3)/4 \Longrightarrow \delta^0(W) \ge \frac{n}{2} + \lambda - 1$.

Disjoint cycles in digraphs

4. Related Results and Open Problems

